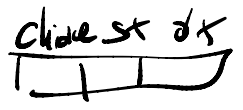


Appendice PUNTORI



```

struct tree_node {
    int key;
    struct tree_node * left;
    struct tree_node * right;
}
    
```

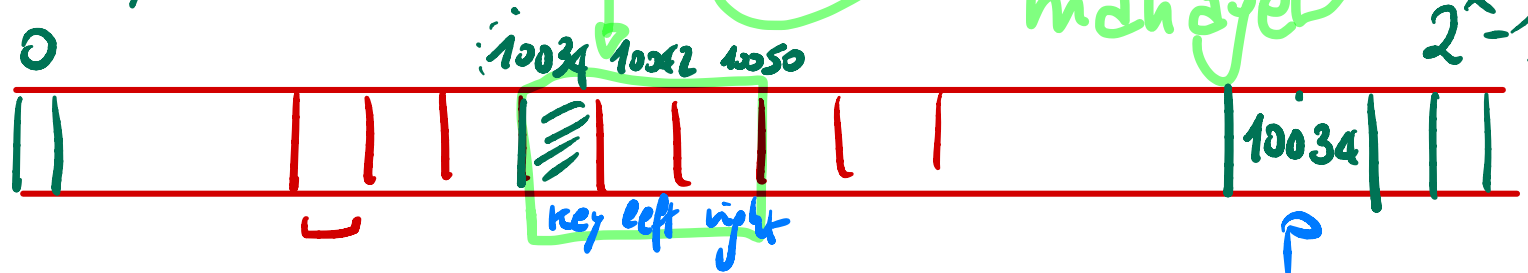
3 celle
di memoria (64 bit (1'uno))

using node =
struct tree_node;

```

auto p = new struct tree_node;
    
```

memory
indirizzo $2^x - 1$



stamp p : 10034

stamp xp :

| | | |
|----|----------------|----------------|
| 32 | i ₁ | i ₂ |
|----|----------------|----------------|

~~10034~~ 10042 10050

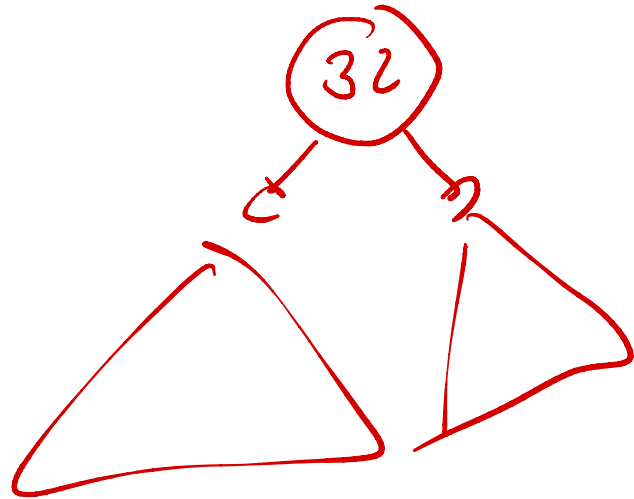
key left right

*p.key ≡ p → key

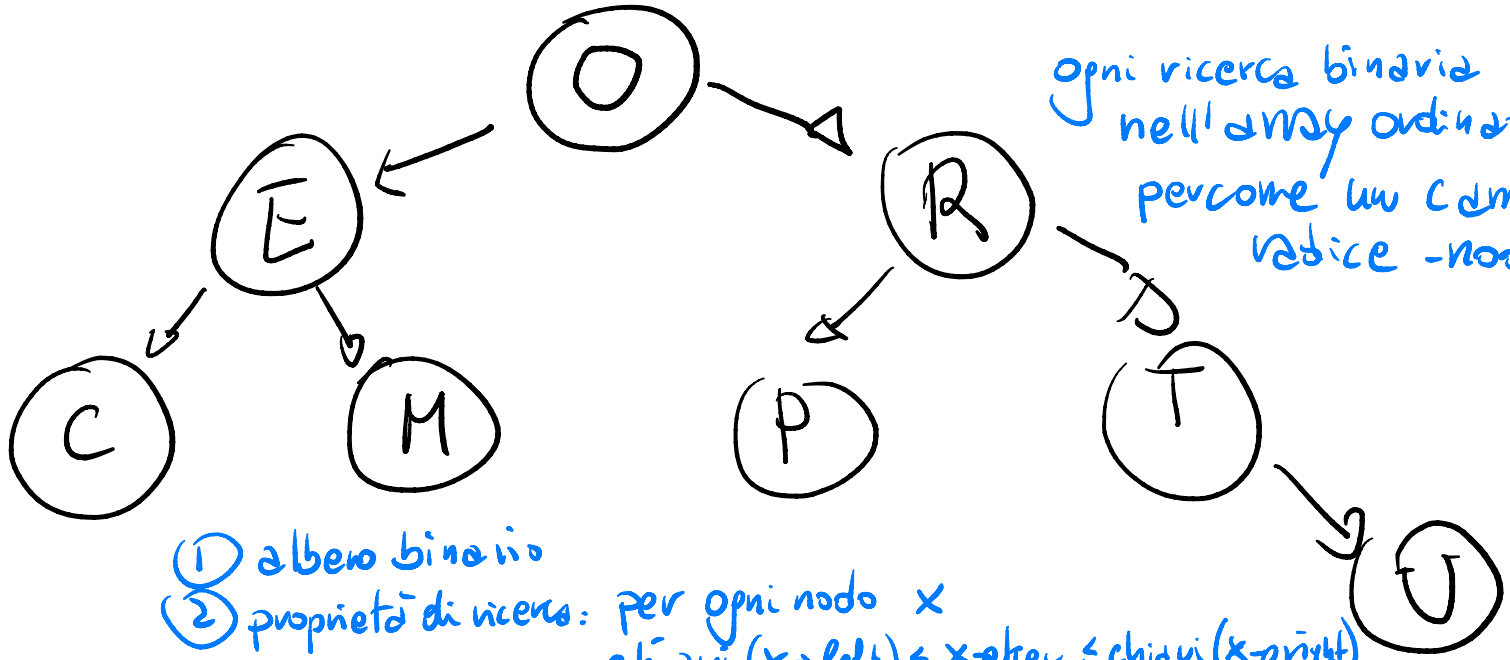
stamp p → key : 32

stamp p → left : i₁

stamp p → right : i₂



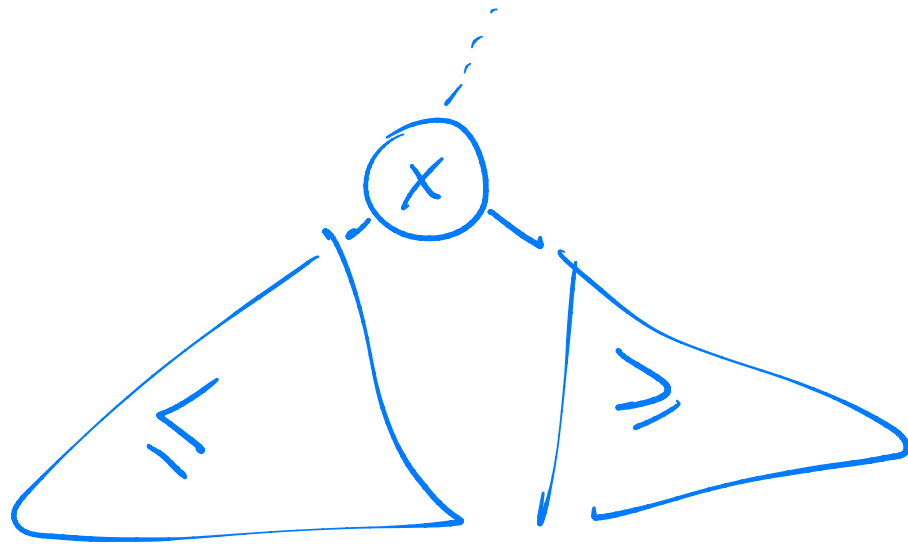
Alberi binari di ricerca



ogni ricerca binaria
nell'array ordinato
percorre un cammino
radice-nodo

- 1) albero binario
- 2) proprietà di ricerca: per ogni nodo x
 $chiar_i(x \rightarrow left) \leq x \rightarrow key \leq chiar_i(x \rightarrow right)$

| 0 | 1 | 2 | 3 | 4 | 5 | 6 | 7 |
|---|---|---|---|---|---|---|---|
| C | E | M | O | P | R | T | U |



struct tree_node * buildTree (vector<int> &A, int sx, int dx)

if (sx > dx) return NULL

auto p = new struct tree_node;

auto cx = random (sx, dx)

// auto cx = $\frac{sx+dx}{2}$

p->key = A[cx]

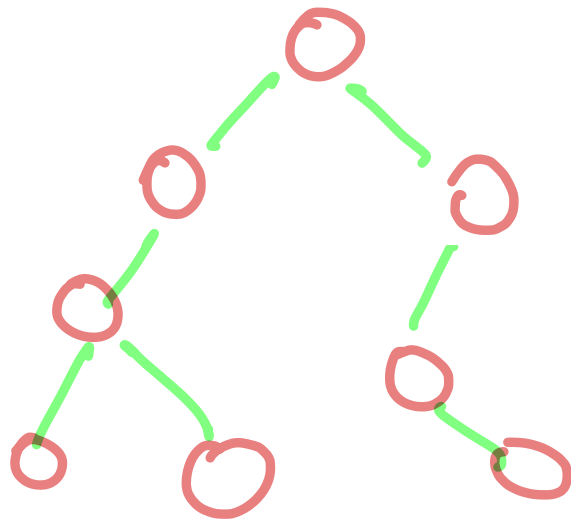
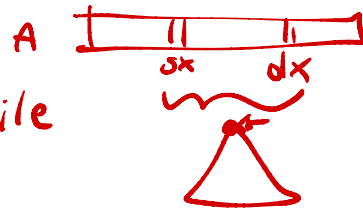
p->left = buildTree (A, sx, cx-1)

p->right = buildTree (A, cx+1, dx)

return p

ORDINATO

decompuibile



RICERCA IN ALBERO BINARIO DI RICERCA

radice
↙
↘ chiave

```
1 Ricerca( u, k ):
2   IF (u == null) RETURN null;
3   IF (k == u.dato.chiave) {
4     RETURN u;
5   } ELSE IF (k < u.dato.chiave) {
6     RETURN Ricerca( u.sx, k );
7   } ELSE {
8     RETURN Ricerca( u.dx, k );
9   }
```

nodo che contiene k
(senza duplicati)

u → key

u → left

u → right

```

1 Ricerca( u, k ):
2   IF (u == null) RETURN null;
3   IF (k == u.dato.chiave) {
4     RETURN u.dato;
5   } ELSE IF (k < u.dato.chiave) {
6     RETURN Ricerca( u.sx, k );
7   } ELSE {
8     RETURN Ricerca( u.dx, k );
9   }

```

ricorsive

```

1 Inserisci( u, e ):
2   IF (u == null) {
3     u = NuovoNodo();
4     u.dato = e;
5     u.sx = u.dx = null;
6   } ELSE IF (e.chiave < u.dato.chiave) {
7     u.sx = Inserisci( u.sx, e );
8   } ELSE IF (e.chiave > u.dato.chiave) {
9     u.dx = Inserisci( u.dx, e );
10  }
11  RETURN u;

```

(post: se k appare già in u, non viene memorizzata)

→ restituire la radice dell'albero esteso con la nuova foglia contenente e

u = new struct tree_node

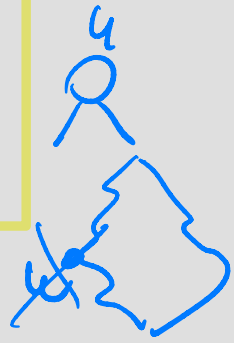
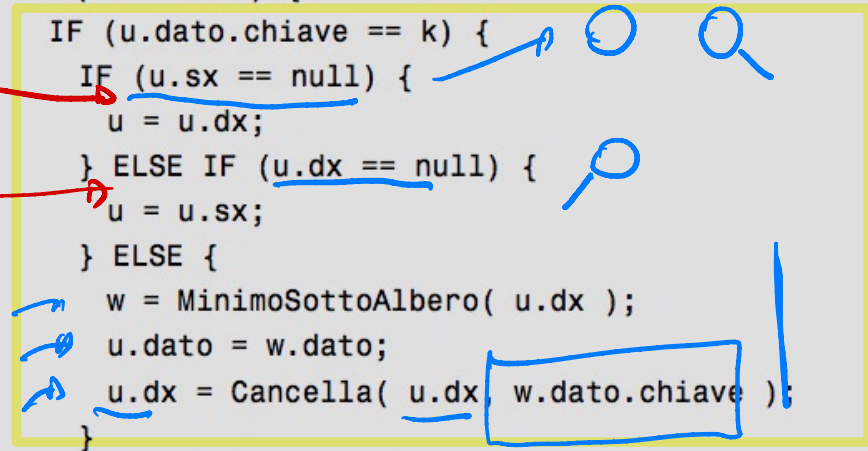
e null null

foglia

```
1 Cancellata( u, k ):
2   IF ( u != null ) {
3     IF ( u.dato.chiave == k ) {
4       IF ( u.sx == null ) {
5         u = u.dx;
6       } ELSE IF ( u.dx == null ) {
7         u = u.sx;
8       } ELSE {
9         w = MinimoSottoAlbero( u.dx );
10        u.dato = w.dato;
11        u.dx = Cancellata( u.dx, w.dato.chiave );
12      }
13    } ELSE IF ( k < u.dato.chiave ) {
14      u.sx = Cancellata( u.sx, k );
15    } ELSE IF ( k > u.dato.chiave ) {
16      u.dx = Cancellata( u.dx, k );
17    }
18  }
19  RETURN u;
```

tmp = u
u = u.dx
delete tmp;

tmp = u;
u = u.sx
delete tmp;



```
1 MinimoSottoAlbero( u ):
2   WHILE ( u.sx != null )
3     u = u.sx;
4   RETURN u;
```

<pre: u ≠ null>

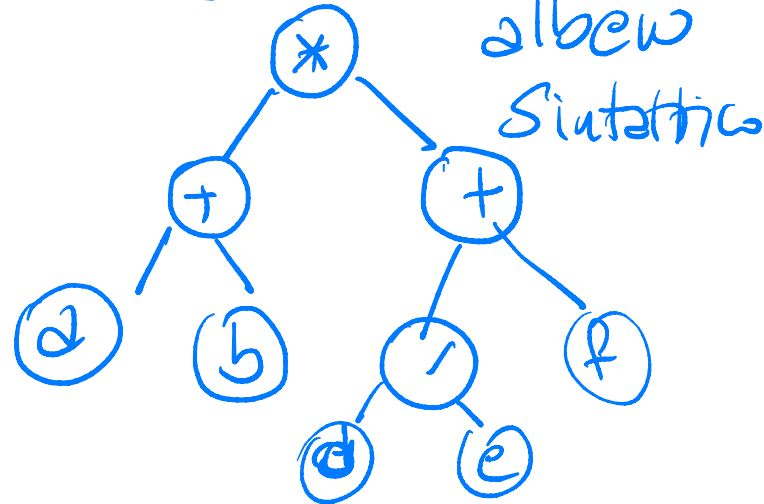

```

void tree-delete (struct tree-node * u) {
    if (u != NULL) {
        tree-delete (u -> left)
        tree-delete (u -> right)
        delete u
    }
}

```

// visits post-order
participate

$$(a+b) * (d/e) + f$$



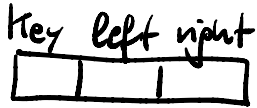
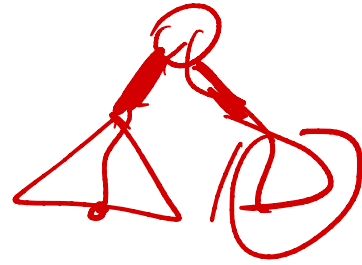
- altezza di ogni nodo

```
- int alt ( u ) {
```

```
    if ( u == NULL ) return -1
```

```
    return 1 + max( alt(u->left), alt(u->right) )
```

```
}
```



```
struct tree_node {
```

```
    int key;
```

```
    int height;
```

```
    struct tree_node * left;
```

```
    struct tree_node * right;
```

```
}
```

invariante

$\forall u: u \rightarrow \text{height} = \text{alt}(u)$

senza eseguire $\text{alt}(u)$!

```

1 Inserisci( u, e ):
2   IF (u == null) {
3     u = NuovoNodo();
4     u.dato = e;
5     u.sx = u.dx = null;
6   } ELSE IF (e.chiave < u.dato.chiave) {
7     u.sx = Inserisci( u.sx, e );
8   } ELSE IF (e.chiave > u.dato.chiave) {
9     u.dx = Inserisci( u.dx, e );
10  }
11  RETURN u;

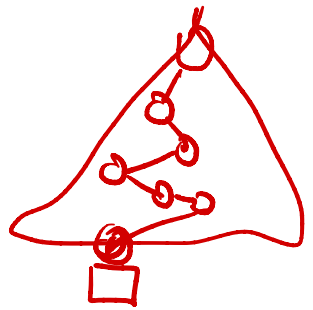
```

$u = \text{new struct tree-node}$

$u \rightarrow \text{height} = 0$

per induzione è quella corretta

$u \rightarrow \text{height} = 1 + \max(\text{myAlt}(u.sx), \text{myAlt}(u.dx))$
 (post: se k appare già in u, non viene memorizzata)



myAlt(u):
 if (u == null) return -1
 return u → height } $O(1)$ tempo

OBIETTIVO : pagare $O(\text{altezza})$ tempo, non $O(\text{n.ro totale di nodi})$!
 mantenere l'altezza durante l'inserimento, evitando di lanciare la funzione alt() perché richiede tempo lineare, altrimenti ...

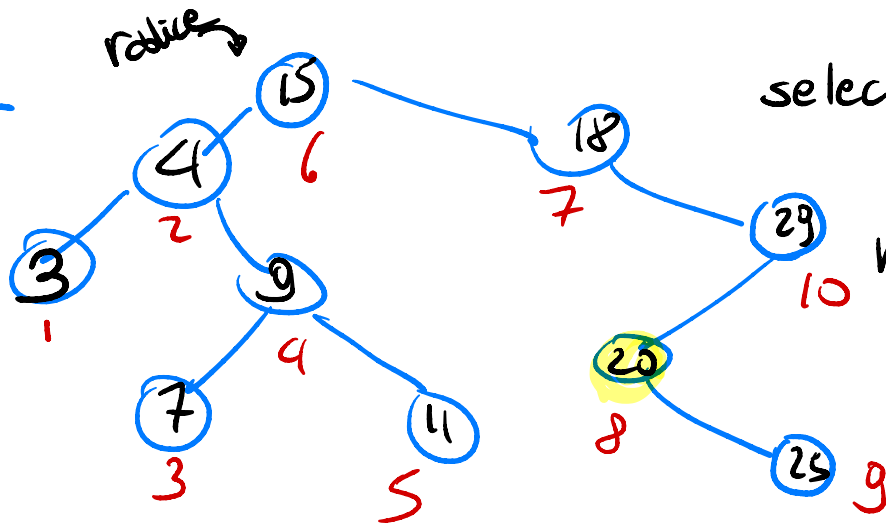
① modificare le insert in modo da mantenere il nuovo campo $u \rightarrow size$

② Utilizzando il campo $u \rightarrow size$

- implementare $rank(u, val) = \# \text{ nodi che hanno } key \leq val$
- " " $select(u, j) = \text{ nodo contenente la } j\text{-esima (de esista)}$

chiave in ordine crescente

esempio



$select(radice, 8) = 20$

$rank(radice, 12) = 5$